Prof. Dr. Andreas Podelski
Dr. Matthias Heizmann

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Discussion: May 14, 2018

Tutorial for Program Verification
Exercise Sheet 3 – Part 1/2

Exercise 1: Hoare logic 2 Points

In this exercise we consider very simple Hoare triples over Boolean variables where

- the precondition *precond*`(X1,...,Xn)` is a Boolean expression over the Boolean variables `X1,...,Xn` and does not contain the Boolean variable `Y`,
- the program consists of the single line

  `Y:=` *expr*`(X1,...,Xn)`,

  where `Y` is a Boolean variable and *expr*`(X1,...,Xn)` is a Boolean expression over the Boolean variables `X1,...,Xn` that does not contain `Y`, and
- the postcondition *postcond*`(X1,...,Xn)` is a Boolean expression over the variables `Y,X1,...,Xn`.

(a) State a propositional logical formula

$$vc(\texttt{Y,X1,...,Xn})$$

that is valid if and only if a Hoare triple that has the following form is valid.

$$\{\ \textit{precond}\texttt{(X1,...,Xn)}\ \}\ \texttt{Y := }\textit{expr}\texttt{(X1,...,Xn)}\ \{\ \textit{postcond}\texttt{(Y,X1,...,Xn)}\ \}$$

(b) Compute your propositional logical formula $vc(\texttt{Z,U,V})$ for the following concrete program.

$$\{\ \texttt{U} \leftrightarrow \texttt{V}\ \}\ \texttt{Z:= U} \wedge \texttt{V}\ \{\ \texttt{Z} \leftrightarrow \texttt{U}\ \}$$

Is your formula valid?

(c) Now we drop the restriction that *precond*`(X1,...,Xn)` does not contain the Boolean variable `Y`. Find a Hoare triple that is not valid but where your formula $vc(\texttt{U,V,Z})$ is valid.

Exercise 2: Hoare logic derivation 2 Points

(a) Write down a partial correctness specification (i.e., precondition and postcondition) for a program `C` that computes the maximum of x and y and stores the result in z.

(b) Write down the program `C`. Use the syntax for programs introduced in the lecture.

(c) Construct a Hoare logic derivation that proves that your program `C` fulfills your correctness specification.

Exercise 3: Hoare triples 2 Points

Consider the following Hoare triples. Which of them are valid for any program `C` and any state assertion ϕ?

(a) $\{\ \mathit{true}\ \}$ `C` $\{\ \phi\ \}$

(b) $\{\ \mathit{false}\ \}$ `C` $\{\ \phi\ \}$

(c) $\{\ \phi\ \}$ `C` $\{\ \mathit{true}\ \}$

(d) $\{\ \phi\ \}$ `C` $\{\ \mathit{false}\ \}$

If a Hoare triple is valid for any program `C` and any state assertion ϕ, then explain why. If a Hoare triple is not valid for some program `C` and some state assertion ϕ, then give a counterexample.