

Termination Analysis by Learning Terminating Programs

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Checking Termination of Programs

- ▶ classical approach

compose termination arguments

- ▶ our approach

decompose program into modules

Checking Termination of Programs

- ▶ classical approach

compose termination arguments

- ▶ our approach

decompose program into modules

Question:

- ▶ What kind of module is suitable for termination proof?

Example: Bubble Sort

```
program sort(int i, int a[])  
   $l_1$  while (i>0)  
     $l_2$    int j:=1  
     $l_3$    while(j<i)  
        if (a[j]>a[i])  
            swap(a,i,j)  
     $l_4$    j++  
   $l_5$    i--
```

Example: Bubble Sort

```
program sort(int i)
```

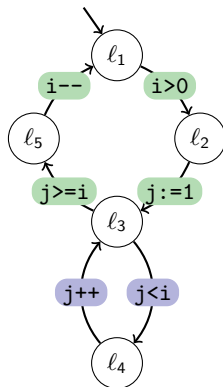
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l2     int j:=1
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l3     while(j<i)
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l4         j++
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l5     i--
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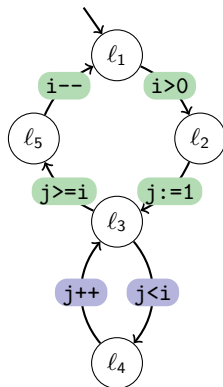
```
l5     i--
```

quadratic ranking function:

$$f(i, j) = i^2 - j$$

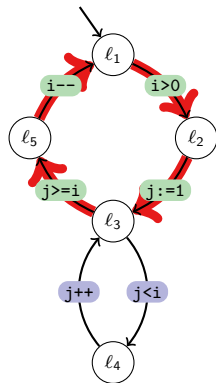
lexicographic ranking function:

$$f(i, j) = (i, i - j)$$



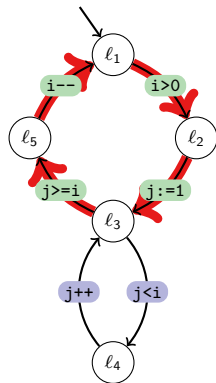
Example: Bubble Sort

single trace OUTER^ω



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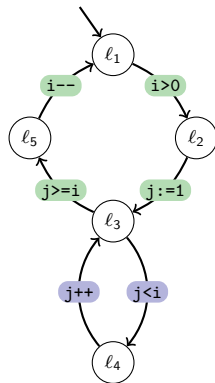
single trace OUTER^ω
has ranking function $f(i, j) = i$



Example: Bubble Sort

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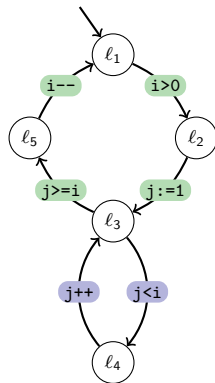
is also ranking function for
set of traces $(\text{INNER}^*.\text{OUTER})^\omega$



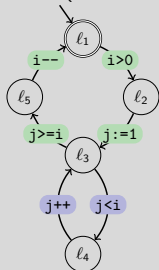
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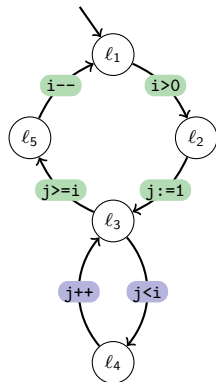
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module \mathcal{P}_1 :
program with fairness constraint whose
set of traces is $(\text{INNER}^*.\text{OUTER})^\omega$

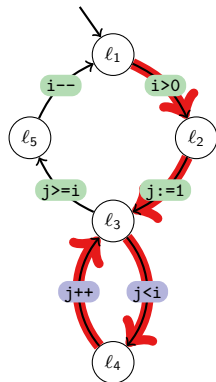


Example: Bubble Sort



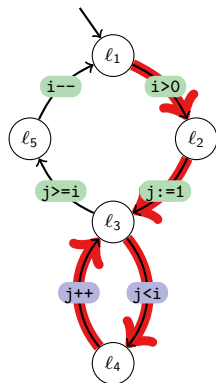
Example: Bubble Sort

new trace OUTER.INNER^ω



Example: Bubble Sort

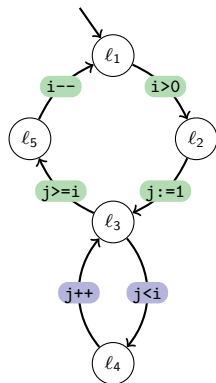
new trace $\text{OUTER.INNER}^\omega$
has ranking function $f(i, j) = i - j$



Example: Bubble Sort

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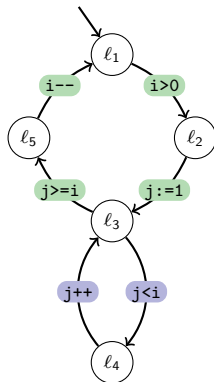
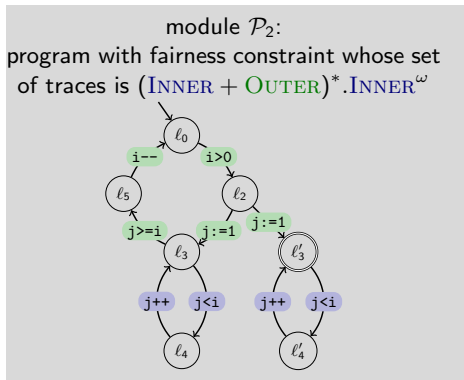
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program \mathcal{P}

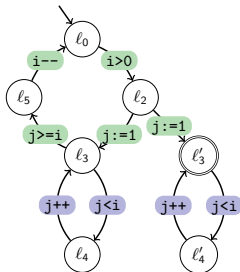
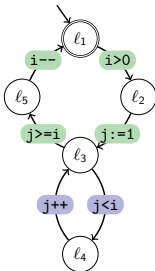
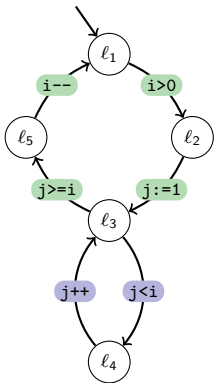
module \mathcal{P}_1

module \mathcal{P}_2

$(\text{OUTER} + \text{INNER})^\omega$

$(\text{INNER}^* \cdot \text{OUTER})^\omega$

$(\text{INNER} + \text{OUTER})^* \cdot \text{INNER}^\omega$



ranking function
 $f(i, j) = i$

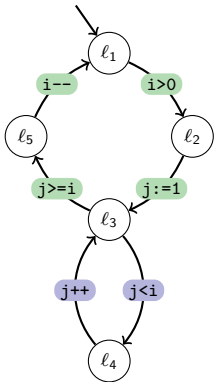
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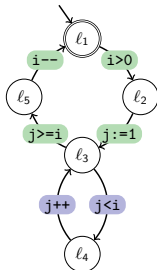
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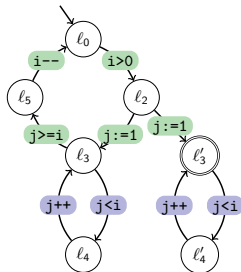
$$(\text{OUTER} + \text{INNER})^\omega = (\text{INNER}^* . \text{OUTER})^\omega + (\text{INNER} + \text{OUTER})^* . \text{INNER}^\omega$$



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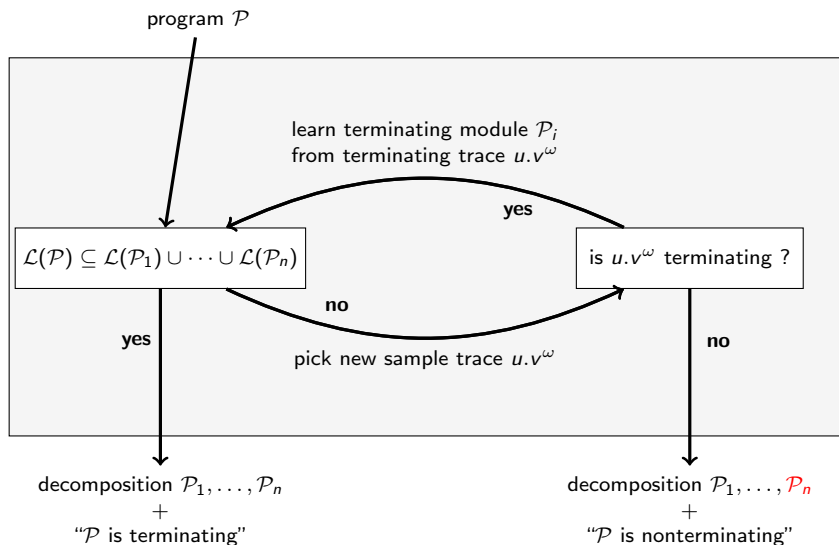
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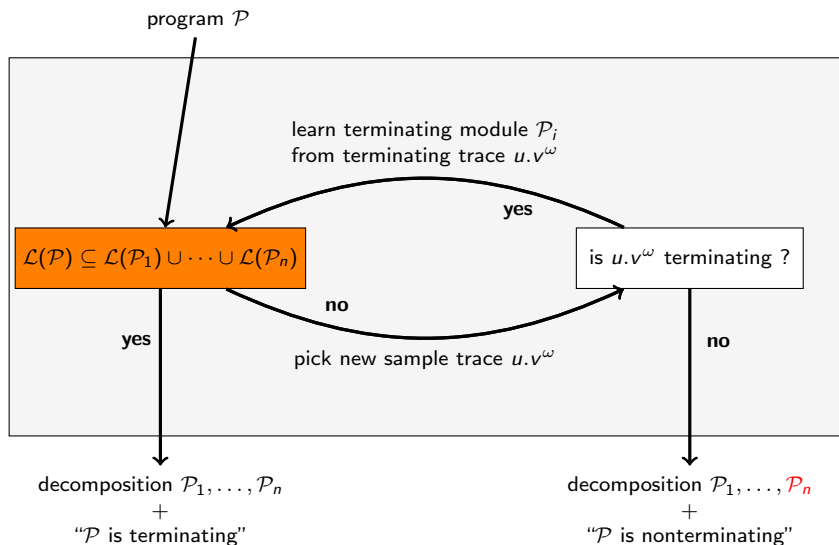
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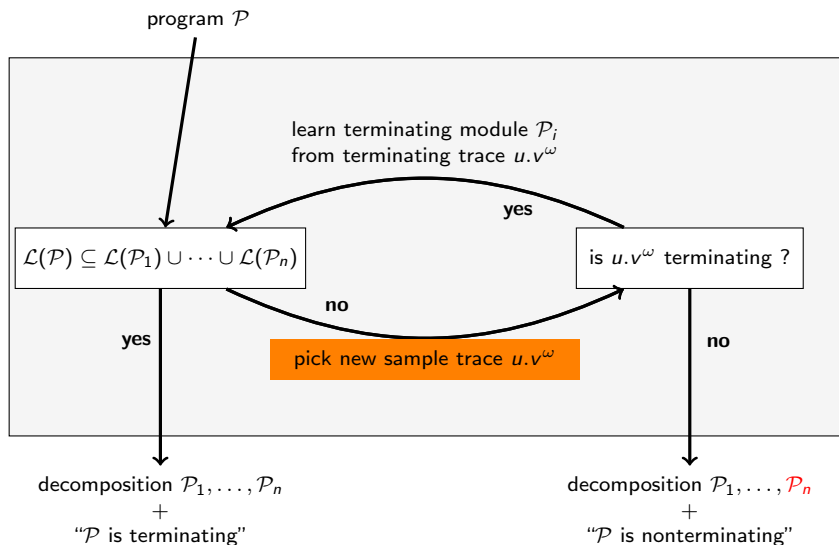
Algorithm for Construction of Decomposition $\mathcal{P}_1, \dots, \mathcal{P}_n$



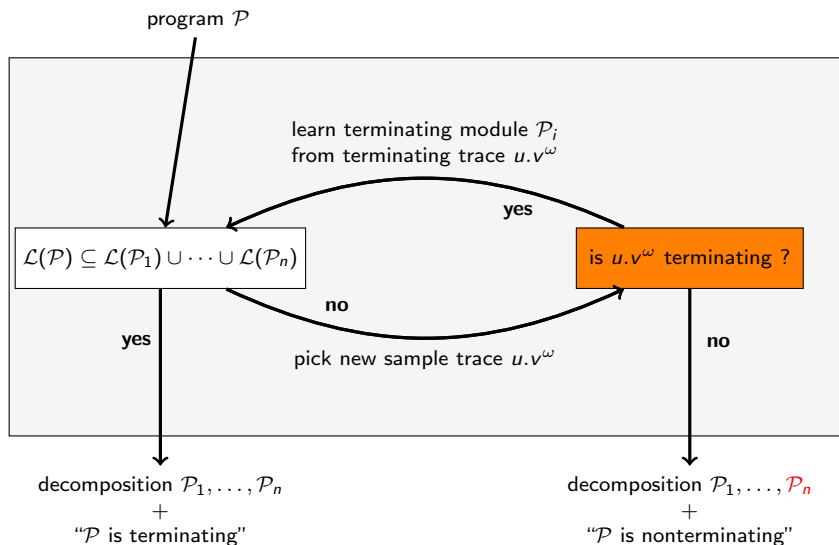
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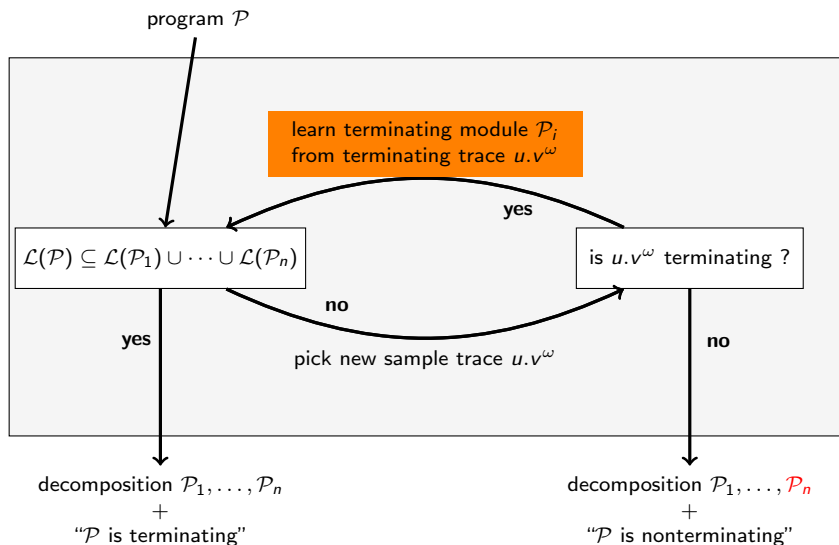
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Learn Module from Trace – Example

input: ultimately periodic trace

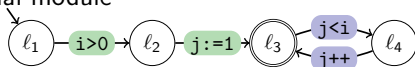
$i>0$ $j:=1$ $(j<i$ $j++)^\omega$,

Learn Module from Trace – Example

input: ultimately periodic trace

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1. construct trivial module

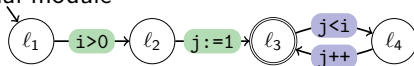


Learn Module from Trace – Example

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2. synthesize ranking function

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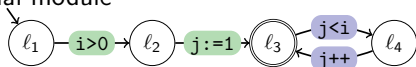
<i>Colón, Sipma</i>	Synthesis of Linear Ranking Functions	(TACAS 2001)
<i>Podelski, Rybalchenko</i>	A complete method for the synthesis of linear ranking functions	(VMCAI 2004)
<i>Bradley, Manna, Sipma</i>	Termination Analysis of Integer Linear Loops	(CONCUR 2005)
<i>Bradley, Manna, Sipma</i>	Linear ranking with reachability	(CAV 2005)
<i>Bradley, Manna, Sipma</i>	The polyranking principle	(ICALP 2005)
<i>Ben-Amram, Genaim</i>	Ranking functions for linear-constraint loops	(POPL 2013)
<i>H., Hoenicke, Leike, Podelski</i>	Linear Ranking for Linear Lasso Programs	(ATVA 2013)
<i>Cook, Kroening, Rümmer, Wintersteiger</i>	Ranking function synthesis for bit-vector relations	(FMSD 2013)
<i>Leike, H.</i>	Ranking Templates for Linear Loops	(TACAS 2014)

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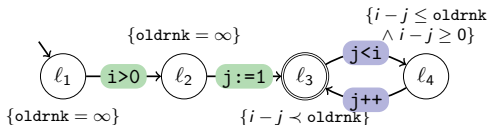
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3. compute rank certificate

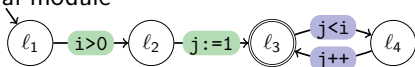


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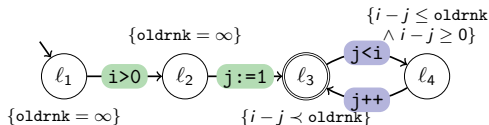
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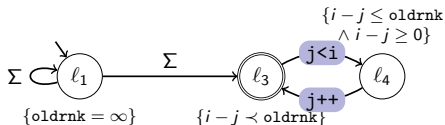
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4. add additional transitions



Our tool:

Ultimate Büchi Automizer

<http://ultimate.informatik.uni-freiburg.de/BuchiAutomizer/>

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Programs with procedures and recursion? Büchi Nested Word Automata!

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1. in our paper: evaluation on benchmark set from

Brockschmidt, Cook, Fuhs **Better termination proving through cooperation** (CAV 2013)

filename	program size	overall runtime	lasso analy time	module co time	Büchi in ch time	modules trivial rf	modules non-trivial	module siz (maximum)
a.10.c.t2.c	183	9s	2.8s	0.7s	2.1s	2	9	5
bF20.t2.c	156	6s	0.7s	0.9s	1.9s	6	7	9
bubbleSort.t2.c	109	5s	0.7s	0.3s	1.2s	5	5	5
consts1.t2.c	40	2s	0.3s	0.1s	0.2s	2	1	5
edn.t2.c	294	119s	18.8s	7.7s	89.0s	141	15	58
eric.t2.c	53	10s	1.1s	1.7s	5.0s	4	6	14
firewire.t2.c	178	28s	3.6s	1.3s	19.0s	12	7	8
mc01.t2.c	47	12s	1.2s	0.6s	4.2s	4	10	8

2. demonstration category on termination in SV-COMP 2014
3. category *C programs* in Termination Competition 2014

Future Work

- ▶ Translate our termination proof (decomposition in terminating modules) into other termination proofs (global ranking function, disjunctive well-founded transition invariant,...) and vice versa.

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- ▶ Translate our termination proof (decomposition in terminating modules) into other termination proofs (global ranking function, disjunctive well-founded transition invariant,...) and vice versa.
- ▶ Find the “simplest” termination argument for a trace $u.v^\omega$.
- ▶ Evaluation of Büchi complementation algorithms in our setting.

Conclusion

- ▶ new termination analysis
- ▶ decompose program \mathcal{P} into modules $\mathcal{P}_1, \dots, \mathcal{P}_n$
- ▶ decomposition not guided by the syntax of the program
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