Ranking Templates for Linear Loops

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Termination

- safety reduced to reachability - liveness reduced to termination

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- safety reduced to reachability - liveness reduced to termination
- neither provable nor refutable by testing
- computing fixpoint on sets of states does not work
- ranking function (decreasing, bounded, contradiction!)

Research directions

1. practical tools for termination analysis

Urban, Miné	**An Abstract Domain to Infer Ordinal-Valued Ranking Functions**	(ESOP 2014)
Brockschmidt, Cook, Fuhs	**Better Termination Proving through Cooperation**	(CAV 2013)
Kroening, Sharygina, Tsitovich, Wintersteiger	**Termination analysis with compositional transition invariants**	(CAV 2010)
Cook, B., Podelski, A., Rybalchenko, A.	**Terminator: Beyond safety**	(CAV 2006)

...

2. decidability of terminaton for restricted classes of programs

Ben-Amram, Genaim	**Ranking functions for linear-constraint loops**	(POPL 2013)
Ben-Amram, Genaim, Masud	**On the Termination of Integer Loops**	(VMCAI 2012)
Tiwari	**Termination of Linear Programs**	(CAV 2004)

...

3. constraint-based synthesis of ranking functions for loops

Cook, Kroening, Rümmer, Wintersteiger	**Ranking function synthesis for bit-vector relations**	(FMSD 2013)
Rybalchenko	**Constraint solving for program verification theory and practice by example**	(CAV 2010)
Colón, Sankaranarayanan, Sipma	**Linear invariant generation using non-linear constraint solving**	(CAV 2003)

...

Ranking functions for loops - applications

- termination analysis for programs
  - Terminator (Cook, Rybalchenko, et al.)
  - T2 (Brockschmidt, et al.)
  - Tan (Chen, Kroening, Wintersteiger, et al.)
  - Ultimate Büchi Automizer (H. et al.)

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- termination analysis for programs
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  - Tan (Chen, Kroening, Wintersteiger, et al.)
  - Ultimate Büchi Automizer (H. et al.)
- cost analysis
- stability of hybrid systems

- affine-linear ranking functions

Colón, Sipma	**Synthesis of Linear Ranking Functions**	(TACAS 2001)
Podelski, Rybalchenko	**A complete method for the synthesis of linear ranking functions**	(VMCAI 2004)
Bradley, Manna, Sipma	**Termination Analysis of Integer Linear Loops**	(CONCUR 2005)
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- multiphase ranking functions

  Bradley, Manna, Sipma **The polyranking principle** (ICALP 2005)

one method to synthesize them all

Ranking function

$Loop(\mathbf{x}, \mathbf{x'})$

Ranking function

$$\forall \mathbf{x}\mathbf{x}'.\ \mathit{Loop}(\mathbf{x}, \mathbf{x}') \rightarrow \underbrace{\mathbf{f}(\mathbf{x}) > \mathbf{f}(\mathbf{x}')}_{\text{decreasing}} \wedge \underbrace{\mathbf{f}(\mathbf{x}) > 0}_{\text{bounded}}$$

Synthesis of ranking function

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Idea:

- write definition as logical formula,
- let theorem prover find satisfying assignment for **free variables**

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- no theorem prover for domain of functions

Solution:

- use template $\mathrm{T}(\mathbf{x}, \mathbf{x}')$

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Synthesis of affine-linear ranking function

$\forall \mathbf{x}\mathbf{x}'.\ \mathit{Loop}(\mathbf{x}, \mathbf{x}') \rightarrow \mathrm{T}(\mathbf{x}, \mathbf{x}')$

where the template $\mathrm{T}(\mathbf{x}, \mathbf{x}')$ is $f(\mathbf{x}) > f(\mathbf{x}') \wedge f(\mathbf{x}) > 0$

decreasing

bounded

and $f(\mathbf{x})$ is a shorthand for the affine-linear term $\mathbf{c}_1 \cdot \mathbf{x}_1 + \cdots + \mathbf{c}_n \cdot \mathbf{x}_1 + \mathbf{c}_0$

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- universal quantification ($\forall \mathbf{x}$ ...)
- nonlinear arithmetic ($\mathbf{c}_1 \cdot \mathbf{x}_1$ ...)

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Lemma (Farkas)

$$\forall x.\ (\ldots \to \ldots) \qquad \mathit{iff} \qquad \exists \vec{\lambda}\ (\ldots)$$

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Lexicographic ranking function

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$$\text{program state} \mapsto \text{lexicographic ordered tuple}$$

Lexicographic order: e.g. $(2, 3, 4) \geq (2, 2, 9)$

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Linear lexicographic ranking function

each entry of the tuple defined by linear function

$$\mathbf{f}(x) = (\ \mathbf{f}_1(x),\ \ldots,\ \mathbf{f}_k(x)\ )$$

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Recall:

Idea:

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Linear lexicographic ranking functions

$$\mathrm{T}(\mathbf{x}, \mathbf{x}') := \mathbf{f}_1(\mathbf{x}) \geq 0 \wedge \mathbf{f}_2(\mathbf{x}') \geq 0 \wedge \left(\mathbf{f}_1(\mathbf{x}) > \mathbf{f}_1(\mathbf{x}') \vee \mathbf{f}_1(\mathbf{x}) \geq \mathbf{f}_1(\mathbf{x}') \wedge \mathbf{f}_2(\mathbf{x}) > \mathbf{f}_2(\mathbf{x}') \right)$$

$\mathbf{f}_1$ bounded

$\mathbf{f}_2$ bounded

$\mathbf{f}_1$ decreasing

$\mathbf{f}_1$ not increasing

$\mathbf{f}_2$ decreasing

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$$\tau(\mathbf{x}, \mathbf{x}') := \mathbf{f}_1(\mathbf{x}) \geq 0 \wedge \mathbf{f}_2(\mathbf{x}') \geq 0 \wedge \left(\mathbf{f}_1(\mathbf{x}) > \mathbf{f}_1(\mathbf{x}') \vee \mathbf{f}_1(\mathbf{x}) \geq \mathbf{f}_1(\mathbf{x}') \wedge \mathbf{f}_2(\mathbf{x}) > \mathbf{f}_2(\mathbf{x}') \right)$$

$\mathbf{f}_1$ bounded

$\mathbf{f}_2$ bounded

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$$\forall \mathbf{x}\mathbf{x}'.\ \mathit{Loop}(\mathbf{x}, \mathbf{x}') \rightarrow \tau(\mathbf{x}, \mathbf{x}')$$

Linear lexicographic ranking functions

$$\tau(\mathbf{x}, \mathbf{x'}) := \underbrace{\mathbf{f}_1(\mathbf{x}) \geq 0}_{\mathbf{f}_1 \text{ bounded}} \wedge \underbrace{\mathbf{f}_2(\mathbf{x'}) \geq 0}_{\mathbf{f}_2 \text{ bounded}} \wedge \Big(\underbrace{\mathbf{f}_1(\mathbf{x}) > \mathbf{f}_1(\mathbf{x'})}_{\mathbf{f}_1 \text{ decreasing}} \vee \underbrace{\mathbf{f}_1(\mathbf{x}) \geq \mathbf{f}_1(\mathbf{x'})}_{\mathbf{f}_1 \text{ not increasing}} \wedge \underbrace{\mathbf{f}_2(\mathbf{x}) > \mathbf{f}_2(\mathbf{x'})}_{\mathbf{f}_2 \text{ decreasing}} \Big)$$

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Lemma (Farkas)

$$\forall x.\ (\ldots \rightarrow \ldots) \qquad \mathit{iff} \qquad \exists \vec{\lambda}\ (\ldots)$$

Theorem (Motzkin)

$$\forall x.\ \neg(\quad \ldots \leq \ldots \quad \wedge \quad \ldots < \ldots \quad) \qquad \mathit{iff} \qquad \exists \vec{\lambda}\ (\ldots)$$

Ranking Template

$\forall \mathbf{x}\mathbf{x}'.\ \mathit{Loop}(\mathbf{x}, \mathbf{x}') \rightarrow \mathrm{T}(\mathbf{x}, \mathbf{x}')$

- "building blocks" linear functions $f(\mathbf{x}) = \mathbf{c}_1 \cdot \mathbf{x}_1 + \cdots + \mathbf{c}_n \cdot \mathbf{x}_1 + \mathbf{c}_0$

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- boolean combinations of linear inequalities (Motzkin applicable)

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Definition (Linear ranking template)

A linear ranking template $\mathrm{T}(\mathbf{x}, \mathbf{x}')$ is a

- boolean combination whose atoms are of the following form

$$\sum_{\mathbf{f} \in F} \alpha_{\mathbf{f}} \cdot \mathbf{f}(\mathbf{x}) + \beta_{\mathbf{f}} \cdot \mathbf{f}(\mathbf{x}') \ \rhd\ 0\ ,$$

  where each $\mathbf{f}(\mathbf{x})$ is an affine-linear term $\mathbf{c}_1 \cdot \mathbf{x}_1 + \cdots + \mathbf{c}_n \cdot \mathbf{x}_1 + \mathbf{c}_0$, each $\alpha_{\mathbf{f}}$, and $\beta_{\mathbf{f}}$ is a constant and $\rhd \in \{\geq, >\}$.
- such that each instance of $\mathrm{T}(\mathbf{x}, \mathbf{x}')$ defines a well-founded relation.

- affine-linear ranking function
  template $\mathrm{T}_{\text{affine}}(\mathbf{x}, \mathbf{x'})$

Lemma

the template $\mathrm{T}_{\text{affine}}(\mathbf{x}, \mathbf{x'})$ is a linear ranking template.

- linear lexicographic ranking function
  k-lexicographic template $\mathrm{T}_{\text{k-lex}}(\mathbf{x}, \mathbf{x'})$

Lemma

the template $\mathrm{T}_{k\text{-}lex}(\mathbf{x}, \mathbf{x'})$ is a linear ranking template, for each k

- piecewise linear ranking function
  k-piece template $\mathrm{T}_{\text{k-piece}}(\mathbf{x}, \mathbf{x'})$

Lemma

the template $\mathrm{T}_{k\text{-}piece}(\mathbf{x}, \mathbf{x'})$ is a linear ranking template, for each k

- multiphase linear ranking function
  k-phase template $\mathrm{T}_{\text{k-phase}}(\mathbf{x}, \mathbf{x'})$

Lemma

the template $\mathrm{T}_{k\text{-}phase}(\mathbf{x}, \mathbf{x'})$ is a linear ranking template, for each k

Why has no one used this method before?

Our explanation: recent progress in solving nonlinear arithmetic

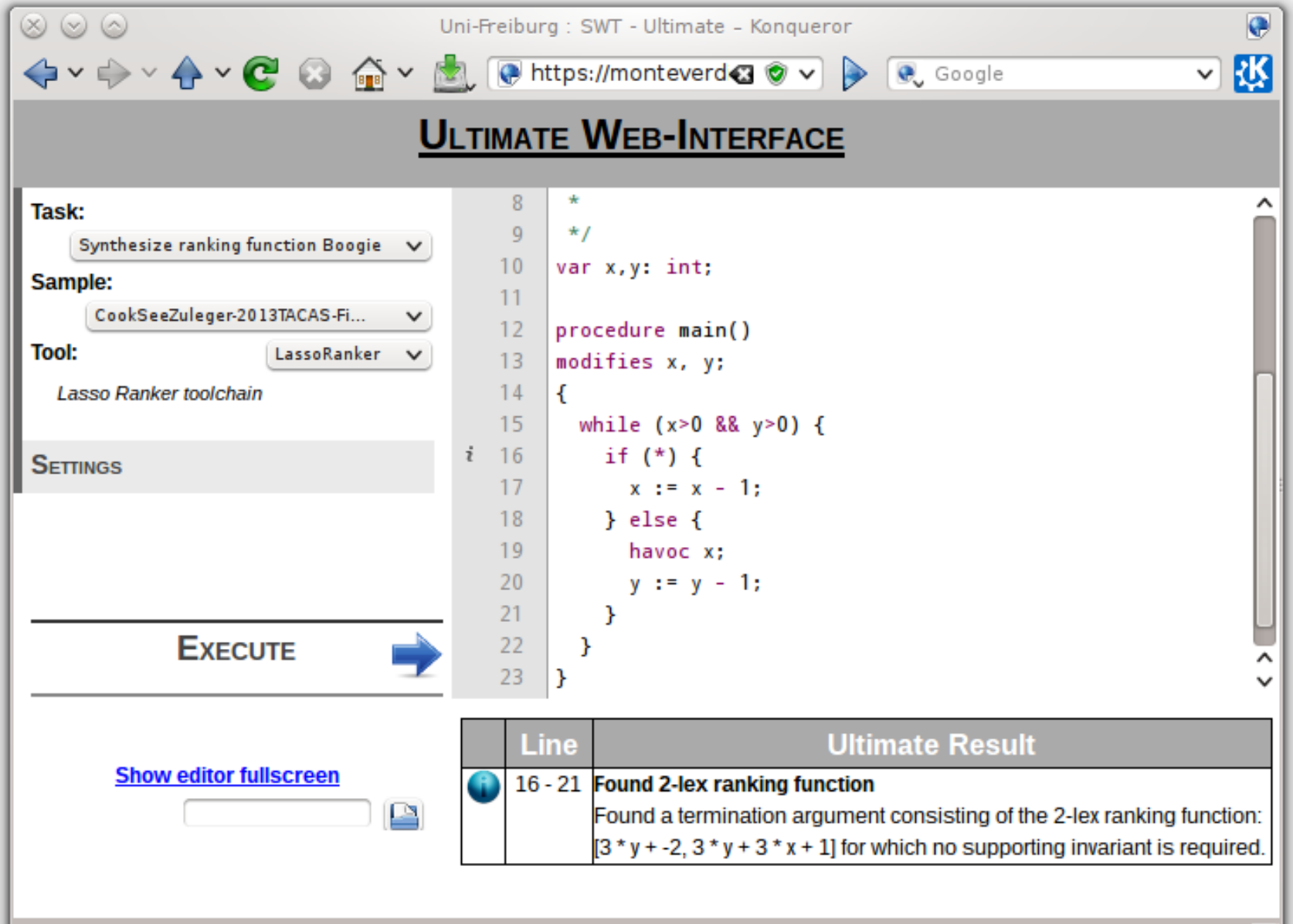
Jovanovic,Moura **Solving non-linear arithmetic** (IJCAR 2012)

SMT solver Z3

`http://z3.codeplex.com/`

Our tool: LassoRanker

`http://ultimate.informatik.uni-freiburg.de/LassoRanker/`



New kind of ranking function?

You can synthesize your new ranking function automatically in three steps.

- write down a template $\mathrm{T}(\mathbf{x}, \mathbf{x'})$ for this ranking function
- prove that each instance of $\mathrm{T}(\mathbf{x}, \mathbf{x'})$ is a well-founded relation
- add template $\mathrm{T}(\mathbf{x}, \mathbf{x'})$ to our tool LassoRanker