

# Generation of Verification Conditions (cont'd)

Andreas Podelski

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- ▶ derivation *unique*
- ▶ verification condition = set of side conditions

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where
$$\begin{aligned}\phi_1 &= \text{wp}(C_1, \psi) \\ \phi_2 &= \text{wp}(C_2, \psi)\end{aligned}$$
- ▶  $\text{wp}(\mathbf{while } b \mathbf{ do } \{ \theta \} C_0, \psi) = \theta$

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- ▶ for command  $C$  of form: **while**  $b$  **do**  $\{\theta\} C_0$  ,
- ▶ add two implications:

$$\phi \rightarrow \theta$$

$$\theta \wedge \neg b \rightarrow \psi$$

and add verification condition for Hoare triple  $\{\theta \wedge b\} C_0 \{\theta\}$

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$$\Gamma \models \Phi \text{ iff } \Gamma \vdash \{\phi\} C \{\psi\}$$